# Simulating with uncertainty : the rough surface scattering problem 

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## A problem often encountered in numerical analysis

(1) A simulation, $\sigma(\mathbf{x}), \mathbf{x} \in \mathbb{D}$ is computationally expensive
© The input x or the domain $\mathbb{D}$ displays uncertainty,
e.g. in simulating the modes of an optical fibre, the refractive index
might not be exactly known, or the boundary may be rough

In this light, more useful than $\sigma(\mathrm{x})$ is the expectation $\langle\sigma(\mathrm{x})\rangle$
Std. dev. in $\sigma(\mathbf{x})$ due to parameter uncertainty also interesting

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## Computing the expectation and standard deviation

## Assumptions

Setup: $d$-dimensional $\mathbf{x}=\left[x_{1}, x_{2}, \ldots, x_{d}\right]$
Each $x_{i}$ is mutually independent, distributed with known pdf $\rho_{i}$ Construct a multi-variate pdf, $\rho(\mathbf{x})=\prod_{i=1}^{n} \rho_{i}\left(x_{i}\right)$

## Quantities of interest



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Quantities of interest

Expectation $\langle\sigma(\mathbf{x})\rangle=\int_{\mathbb{D}} \sigma(\mathbf{x}) \rho(\mathbf{x}) d \mathbf{x}$
Std. dev. $\quad \Delta \sigma(\mathbf{x})=\sqrt{\int_{\mathbb{D}}(\sigma(\mathbf{x})-\langle\sigma(\mathbf{x})\rangle)^{2} \rho(\mathbf{x}) d \mathbf{x}}$

## Any problem with $\langle\sigma(\mathbf{x})\rangle=\int_{\mathbb{D}} \sigma(\mathbf{x}) \rho(\mathbf{x}) d \mathbf{x}$ ?

Yes!
$\sigma(\mathrm{x})$ is only known numerically.
Recall: $\sigma(\mathbf{x})$ is the output of an (possibly expensive) simulation

Monte Carlo method
The most commonly used method to estimate expectation.
Instantiate several $\mathbf{x}_{i}$ 's and approximate $\langle\sigma(\mathbf{x})\rangle=\sum_{i=1}^{n} \frac{1}{n} \sigma\left(\mathbf{x}_{i}\right)$
Convergence rate: independent of $d$, but slow $O\left(\frac{1}{\sqrt{n}}\right)$
Extra: Create a histogram from $\sigma\left(\mathbf{x}_{i}\right)$ values to estimate pdf of $\sigma$

## Objective of this talk: Discuss alternatives to Monte Carlo

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## The alternatives to Monte Carlo

Broadly, two families of methods will be discussed:
(1) Galerkin Polynomial Chaos (gPC)

Need to code a new solver, works for small range of $d$ values
(2) Stochastic Collocation (SC)

Uses an existing solver, works for small-medium range of $d$ values Attractive, because compatible with commercial software

## What about Monte Carlo? <br> Use for benchmarking and for large range of $d$ values

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## Start with Stochastic Collocation: Polynomial interpolation

For simplicity, consider a 1-D case of the simulation, $\sigma(x)$
Polynomial interpolation and integration

- Choose $n$-points to evaluate $\sigma(x)$ and construct a $n-1$ degree polynomial (Lagrange interpolation): $\sigma_{n-1}(x)=\sum_{i=1}^{n} \sigma\left(x_{i}\right) L_{i}(x)$
Recall, $L_{i}(x)=\prod_{j=1, j \neq i}^{n}\left(x-x_{j}\right) /\left(x_{i}-x_{j}\right)$
- Compute expectation as

- $\alpha_{i}$ can be pre-computed, and $\langle\sigma(x)\rangle$ is accurate to order $n-1$.


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$\langle\sigma(x)\rangle \approx \sum_{i=1}^{n} \sigma\left(x_{i}\right) \int L_{i}(x) \rho(x) d x=\sum_{i=1}^{n} \sigma\left(x_{i}\right) \alpha_{i}$
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Can we do better?

## Stochastic Collocation: Gaussian quadrature

Yes! Gaussian quadrature

- Use the theory of orthogonal polynomials, i.e. find polynomials s.t. $\int p_{m}(x) p_{n}(x) \rho(x) d x=\delta_{m, n}$
- Pick the $n$ points, $\left\{x_{i}\right\}$ to be roots of $n^{t h}$ degree polynomial, $p_{n}(x)$
- This also gives $\langle\sigma(x)\rangle \approx \sum_{i=1}^{n} \sigma\left(x_{i}\right) \alpha_{i}$, but now integral is accurate to order $2 n-1$
- Legendre, $x \in[-1,1], \rho(x)=1$
- Jacobi, $x \in(-1,1), \rho(x)=(1-x)^{\alpha}(1+x)^{\beta}, \alpha, \beta>-1$
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Common orthogonal polynomials

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## But our problem is $d$-dimensional!

Solution: Extend one-dimensional Gaussian quadrature to $d$ dimensions

- With $\mathbf{x}=\left[x_{i}, x_{2}, \ldots, x_{d}\right]$, express $\sigma(\mathbf{x})=\prod_{i=1}^{d} \sigma_{i}\left(x_{i}\right)$ (implicitly)
- The integral splits up: $\int \sigma(\mathbf{x}) \rho(\mathbf{x}) d \mathbf{x}=\prod_{i} \int \sigma_{i}\left(x_{i}\right) \rho_{i}\left(x_{i}\right) d x_{i}$
- Apply $n$-point Gaussian quadrature (GQ) in each dimension

- An example in 2D $(x, y)$ and a 2-point GQ in each dimension $\langle\sigma(x, y)\rangle=\left[\sigma_{x}\left(x_{1}\right) \alpha_{1}+\sigma_{x}\left(x_{2}\right) \alpha_{2}\right]\left[\sigma_{y}\left(y_{1}\right) \alpha_{1}+\sigma_{y}\left(y_{2}\right) \alpha_{2}\right]=$ $\sigma\left(x_{1}, y_{1}\right) \alpha_{1,1}+\sigma\left(x_{1}, y_{2}\right) \alpha_{1,2}+\sigma\left(x_{2}, y_{1}\right) \alpha_{2,1}+\sigma\left(x_{2}, y_{2}\right) \alpha_{2,2}$

All dimensions are multiplied: called the tensor product rule

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- Apply $n$-point Gaussian quadrature (GQ) in each dimension: $\langle\sigma(\mathbf{x})\rangle=\prod_{i=1}^{d} \sum_{j=1}^{n} \sigma_{i}\left(x_{i, j}\right) \alpha_{j}$. Combine the $\sigma_{i}$ 's to get $\sigma\left(\mathbf{x}_{k}\right)$.
- An example in 2D $(x, y)$ and a 2-point GQ in each dimension: $\langle\sigma(x, y)\rangle=\left[\sigma_{x}\left(x_{1}\right) \alpha_{1}+\sigma_{x}\left(x_{2}\right) \alpha_{2}\right]\left[\sigma_{y}\left(y_{1}\right) \alpha_{1}+\sigma_{y}\left(y_{2}\right) \alpha_{2}\right]=$ $\sigma\left(x_{1}, y_{1}\right) \alpha_{1.1}+\sigma\left(x_{1}, y_{2}\right) \alpha_{1.2}+\sigma\left(x_{2}, y_{1}\right) \alpha_{2.1}+\sigma\left(x_{2}, y_{2}\right) \alpha_{2.2}$ All dimensions are multiplied: called the tensor product rule


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\begin{aligned}
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& \sigma\left(x_{1}, y_{1}\right) \alpha_{1,1}+\sigma\left(x_{1}, y_{2}\right) \alpha_{1,2}+\sigma\left(x_{2}, y_{1}\right) \alpha_{2,1}+\sigma\left(x_{2}, y_{2}\right) \alpha_{2,2}
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All dimensions are multiplied: called the tensor product rule

## Visualizing the tensor product rule

- E.g. function evaluation points in a 2D 5-point GQ (25 evals) : Denote as $\langle\sigma\rangle_{5,5}$
- Curse of dimensionality is clear: number of function evaluations $=n^{d}$
- Can we do better?

Theorem: (Mysovskikh 1968, Möller 1976)
To attain a polynomial exactness equal to $m$, the (optimal) required number of grid-points has lower and upper bounds given by


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N_{\min }=\binom{d+\lfloor m\rfloor / 2}{\lfloor m\rfloor / 2} \leq N_{\mathrm{opt}} \leq\binom{ d+m}{m}=N_{\max }
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## Sparse Grids: Smolyak (1963)

(1) Every $\langle\sigma\rangle_{i, j}$ is an approximation to the actual integral, $\langle\sigma\rangle$
(2) Introduce a "level" parameter $k=(i+j)$ and let the max level be denoted by $l=\max \{k\}$

## Telescope a series of different levels to approximate $\langle\sigma\rangle$

- e.g. (Max) level $l=4$ :
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As the max level increases, approximation becomes better

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$\langle\sigma\rangle \approx\left[\langle\sigma\rangle_{3,1}+\langle\sigma\rangle_{1,3}+\langle\sigma\rangle_{2,2}\right]_{\mathrm{k}=4}-\left[\langle\sigma\rangle_{2,1}+\langle\sigma\rangle_{1,2}\right]_{\mathrm{k}=3}$
- e.g. (Max) level $l=5$ :
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Sparse Grid Points


## 13 Points (SG) v/s 25 points (TP)!

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Tensor Grid Points

Substantial
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## Sparse Grid rule: finer points

TP rule: $\langle\sigma\rangle=\sum_{i_{1}=1}^{n_{1}} \ldots \sum_{i_{d}=1}^{n_{d}} \sigma\left(x_{i_{1}}, \ldots, x_{i_{d}}\right) \alpha_{i_{1}} \ldots \alpha_{i_{d}}$
$=\left(Q^{n_{1}} \times Q^{n_{2}} \times \ldots \times Q^{n_{d}}\right)[\sigma]$
$n_{i}$ point quadrature in the $i^{\text {th }} \mathrm{dim} ; \prod n_{i} \approx n^{d}$ points


- Nested quadrature rules (e.g. Clenshaw-Curtis, Gauss-Konrad) re-use fn evaluation points between levels.


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## Switching gears: from sampling to projection

The two methods (Monte Carlo, Stochastic Collocation) considered so far were of a "sampling" kind: $\langle\sigma\rangle$ estimated using samples of $\sigma(\mathbf{x})$.


Straightforward when x is spatio-temporal. But when stochastic?

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Projection based approach: Overview of Galerkin method

- Governing equation: $\Theta f(\mathbf{x})=g(\mathbf{x})$, where $\Theta$ is an operator, $g$ is a known function, and $f$ is to be determined.
- Project $f$ in a known basis, $\left\{\phi_{j}\right\}: f(\mathrm{x})=\sum_{j} u_{j} \phi_{j}(\mathrm{x})$
- Take an inner product with the same basis functions on both sides to get: $\sum_{j}\left\langle\phi_{i}(\mathbf{x}), \Theta \phi_{j}(\mathbf{x})\right\rangle u_{j}=\left\langle\phi_{i}(\mathbf{x}), g(\mathbf{x})\right\rangle . B C$ used to simplify.
- This is a system of equations; solve for $u$ and get $f$

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## generalized Polynomial Chaos (gPC): a very brief history

- Polynomial Chaos (PC): coined by Norbert Wiener studying Gaussian stochastic processes (1938). Used Hermite polynomials as basis.
- Ghanem (1998) used theory of Wiener-Hermite PC to represent random processes in an orthogonal basis of Hermite polynomials.
- Xiu, Karniadakis (2002) generalize to non-Gaussian using other orthogonal polynomials, wavelets, etc: generalized PC (gPC)

Finally, solve the gPC system of equations using galerkin projection: stochastic galerkin (SG) method.

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## generalized Polynomial Chaos (gPC): a very brief history

- Polynomial Chaos (PC): coined by Norbert Wiener studying Gaussian stochastic processes (1938). Used Hermite polynomials as basis.
- Ghanem (1998) used theory of Wiener-Hermite PC to represent random processes in an orthogonal basis of Hermite polynomials.
- Xiu, Karniadakis (2002) generalize to non-Gaussian using other orthogonal polynomials, wavelets, etc: generalized PC (gPC)

Finally, solve the gPC system of equations using galerkin projection: stochastic galerkin (SG) method.

## The basis functions in gPC

- Start with a RV, call it $z$ as before. Let it have a distribution function, $F_{z}(\theta)=P(z \leq \theta)$ and a pdf $\rho(\theta)$ s.t. $d F_{z}(\theta)=\rho(\theta) d \theta$
- The generalized Polynomial Chaos basis functions are orthogonal basis functions, $\psi_{i}(z)$, satisfying $\left\langle\psi_{i}(z) \psi_{j}(z)\right\rangle=\int \psi_{i}(\theta) \psi_{j}(\theta) \rho(\theta) d \theta=\gamma_{i} \delta_{i j}$
- Construct linear space of polynomials of degree at most $n: \mathbb{P}_{n}(z)$
- Various kinds depending on $\rho(\theta)$
- Legendre, $\theta \in[-1,1], \rho(\theta)=1 / 2$.
- Jacobi, $\theta \in(-1,1), \rho(\theta)=(1-\theta)^{\alpha}(1+\theta)^{\beta}, \alpha, \beta>-1$.
- Hermite, $\theta \in(-\infty, \infty), \rho(\theta)=e^{-\theta}$


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## A scalar example of gPC-SG

- Consider: a simple equation in one unknown, $u: a u=b$ Now: Let the system parameters $a, b$ have some uncertainty, e.g. $a(\theta)=a_{0}+\alpha \theta$, where $\theta$ is a uniform RV in $[-0.5,0.5]$. In this case what is $\langle u\rangle$ or $\left\langle u^{2}\right\rangle$ ?
- Expand: solution in the $\psi$ basis $-u(\theta)=\sum_{j=1}^{n} u_{j} \psi_{j}(\theta)$
- Do: Galerkin testing with the same basis functions: Get a system of equations, where we solve for $u$ : $A u=b$, where

So: $\langle u\rangle=\int u(\theta) \rho(\theta) d \theta=\sum_{i} u_{i}\left\langle\psi_{i}(\theta)\right\rangle$ and $\left\langle u^{2}\right\rangle=\sum_{i} u_{i}^{2} \gamma_{i}$.
$\Longrightarrow$ std. dev. in $u$ can be computed : $\sqrt{\left\langle u^{2}\right\rangle-\langle u\rangle^{2}}$
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- $s_{0}(x)$ is the mean of the random process
- $\eta, f$ solve this eigenvalue problem: $\int C(i, j) f_{k}(j) d j=\eta_{k} f_{k}(i)$ where
- $z(\theta)$ represents mutually uncorrelated normal $\operatorname{RVs}\left(\left\langle z_{k}\right\rangle=0\right)$
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## KL expansion for exponential correlation

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s(x, \theta)=s_{0}(x)+\sum_{k=1}^{\infty} \sqrt{\eta_{k}} f_{k}(x) z_{k}(\theta)
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exponential correlation function
$C(i, j)=\exp (-|i-j| / l)$,
(correlation length $l$ ).
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## The case of multiple random variables with correlation

Consider again the random rough surface ...

given by the KL expansion: $s(x, \theta)=s_{0}(x)+\sum_{k=1}^{d} \sqrt{\eta_{k}} f_{k}(x) z_{k}(\theta)$ consists of multiple random variables, $z_{i}$.

- Multivariate gPC (i.e. tensor product of univariate gPC ) $\Psi_{\mathbf{i}}(\vec{z})=\psi_{i_{1}}\left(z_{1}\right) \ldots \psi_{i_{d}}\left(z_{d}\right), \quad 0 \leq \sum_{j=1}^{d} i_{j} \leq n$ (max degree) where $\mathrm{i}=\left(i_{1}, \ldots, d\right)$ : indices, $\vec{z}=\left(z_{1}, \ldots, z_{d}\right)$ RVs Belong to the space $\mathbb{P}_{n}^{d}$ of dimension $w=\binom{n+d}{n}$
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## Changing gears ...

So far ...
This completes an overview of stochastic computation:
(1) Monte Carlo (MC)
(2) Stochastic Collocation (SC)
(3) Stochastic Galerkin (SG) using generalized Polynomial Chaos
$\qquad$
$\square$ electromagnetic scattering from a random rough surface: e.g. seen in microwave remote sensing, $[\mathrm{MC}]^{a},[S C, S G]^{b}$
> ${ }^{a}$ Khankhoje et al. "Computation of radar scattering from heterogeneous rough soil using the finite element method", 2013 IEEE TGRS
> ${ }^{b}$ Khankhoje et al. "Stochastic Solutions to Rough Surface Scattering using the finite element method"

> 2017 IEEE TAP

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Moving on ...
To make matters more concrete, consider the problem of computing the electromagnetic scattering from a random rough surface:
e.g. seen in microwave remote sensing, $[\mathrm{MC}]^{a},[\mathrm{SC}, \mathrm{SG}]^{b}$

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## Traditional FEM setup for rough surface scattering



- Random rough surface instance generated, and the domain meshed
- Based on incident field, Radar cross-section (RCS) computed
- Above steps repeated for several $(\approx 100)$ instances
- Quite inefficient!


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Partition the domain into parts that can move, and those that

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\Delta y=\left\{\begin{array}{l}
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## The standard FEM recipe: deterministic solver

2D vector FE basis functions $\vec{W}$


Maxwell's equations in weak form:

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\int \vec{W} \cdot\left[\nabla \times\left(\frac{1}{\epsilon_{r}} \nabla \times \vec{H}\right)-k_{0}^{2} \mu_{r} \vec{H}\right] d S=0
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\begin{aligned}
A \mathbf{u} & =\mathbf{b}, A \in \mathbb{C}^{m \times m}, \mathbf{u}, \mathbf{b} \in \mathbb{C}^{m}, \\
A_{p q} & =\sum_{e} \alpha_{e, p q}\left(\vec{r}_{e}\right)+\delta_{p q} \nu_{p}\left(\vec{r}_{p}\right), \\
b_{p} & =\tau_{p}\left(\vec{r}_{p}\right), \quad \text { where } \\
\vec{r}_{e} & =\left(x_{i}, x_{j}, x_{k}, y_{i}, y_{j}, y_{k}\right), e^{\mathrm{th}} \text { ele } \\
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Quantity of interest: radar cross-section $\sigma_{2 D}=\lim _{r \rightarrow \infty} 2 \pi r\left|\frac{E_{z}^{f}}{E_{z}^{i}}\right|^{2}$
$E_{z}^{f}(\vec{r})=\sqrt{\frac{k_{0}}{8 \pi}} \frac{e^{-i\left(k_{0} r-\pi / 4\right)}}{\sqrt{r}} \times \oint \hat{z} \cdot\left(\hat{r} \times \vec{M}\left(\overrightarrow{r^{\prime}}\right)+Z_{0} \mu_{r} \hat{r} \times \hat{r} \times \vec{J}\left(\overrightarrow{r^{\prime}}\right)\right) e^{i k_{0} \hat{r} \cdot \vec{r}^{\prime}} d l^{\prime}$
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till convergence ( $\approx 100$ times)


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## What is deterministic here?

Input to the FEM solver:
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The FEM is run for a specified surface, hence "deterministic" solver

- Monte Carlo process converges as $\mathbb{O}\left(1 / \sqrt{n_{m c}}\right)$, independent of $d$. (Recall: $d$ depends on the correlation length of the surface) i.e. $\langle\sigma\rangle=\sum_{i=1}^{n_{m c}} \sigma\left(\vec{r}, \vec{z}_{i}\right) / n_{m c}$
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## From Monte Carlo to Stochastic Collocation

(1) Construct multivariate pdf of the $d$ random normal variables: $\rho(\vec{z})=\prod_{j=1}^{d} \rho_{j}\left(z_{j}\right)$ over domain $\mathcal{D}^{d}, \mathcal{D}=(-\infty, \infty)$
(2) Express expected value as $\langle\sigma\rangle=\int_{\mathcal{D}^{d}} \sigma(\vec{r}, \vec{z}) \rho(\vec{z}) d \vec{z}$
(3) Consider $n_{s c}$ evals of the above integral at predecided quadrature points, $\overrightarrow{z_{i}}$.
(1) Express $\sigma(\vec{r}, \vec{z})$ in terms of interpolating multivariate polynomials (e.g. Langrage) $\left\{P^{(i)}(\vec{z})\right\}_{i=1}^{d}$ giving $\sigma(\vec{r}, \vec{z})=\sum_{i=1}^{n_{s c}} \sigma\left(\vec{r}, \overrightarrow{z_{i}}\right) P^{(i)}$
Finally, $\langle\sigma\rangle=\sum_{i=1}^{n_{s c}} \sigma\left(\vec{r}, \vec{z}_{i}\right) \alpha_{i}$ where $\alpha_{i}=\int_{\mathcal{D}^{d}} \rho(\vec{z}) P^{(i)}(\vec{z}) d \vec{z}$

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## Stochastic Collocation for rough surface scattering

SC recipe: $\langle\sigma\rangle=\sum_{i=1}^{n_{\text {sc }}} \sigma\left(\vec{r}, \vec{z}_{i}\right) \alpha_{i}$
This just looks like the old MC formula! $\langle\sigma\rangle=\sum_{i=1}^{n_{m c}} \sigma\left(\vec{r}, \vec{z}_{i}\right) / n_{m c}$
So, we can use the same solver for $\sigma\left(\vec{r}, \vec{z}_{i}\right)$

## Results

How many function
evals using sparse grid? $\approx 2^{l} d^{l} / l$
$\leftarrow$ Compare MC and SC for $d=50, l=1$
Evals:
MC=100, SC=101

## Stochastic Collocation for rough surface scattering

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## Conclusions about Stochastic Collocation

- Given surface needs to be accurately represented by the finite KL expansion, i.e. there is an optimal $d$.
- For MC, found that $\approx 100$ iterations give convergence to within 1dB.
- The cheapest SC would have level $l=1$, giving $n_{s c} \approx 2 d$ $\Longrightarrow$ if the surface requires $d>50, \mathrm{MC}$ is better
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## Moving over to Stochastic Galerkin

- Recall the earlier (scalar) example of $a u=b$, where we replaced $a \rightarrow a_{0}+\alpha \theta$, where $\theta$ was a uniform RV.
- In the case of the FEM, we have a sparse matrix equation $\mathrm{Au}=\mathrm{b}$ to solve. Now, $A_{p q}$ must be transformed as per the KL expansion for the surface. Recall:

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\begin{aligned}
A_{p q} & =\sum_{e} \alpha_{e, p q}\left(\vec{r}_{e}\right)+\delta_{p q} \nu_{p}\left(\vec{r}_{p}\right), \\
b_{p} & =\tau_{p}\left(\vec{r}_{p}\right), \quad \text { where } \\
\vec{r}_{e} & =\left(x_{i}, x_{j}, x_{k}, y_{i}, y_{j}, y_{k}\right), e^{\text {th }} \text { ele } \\
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Use our mesh deformation scheme, which transformed $y \rightarrow y+\Delta y$.
This leads to $A_{p q} \rightarrow \tilde{A}_{p q}, b_{p} \rightarrow \tilde{b}_{p}$

## Steps in Stochastic Galerkin

(1) Accomplish the perturbations by Taylor expanding to second order:
( $1^{\text {st }}, 2^{\text {nd }}$ derivatives are computed by finite differences)

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& \tilde{A}_{p q}=\beta_{p q}+\sum_{i=1}^{d} \beta_{p q}^{(i)} z_{i}+\sum_{i, j=1}^{d} \beta_{p q}^{(i, j)} z_{i} z_{j}, \leftarrow \text { matrices } m \times m \\
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(3) Now do Galerkin testing in the basis of orthogonal polynomials corresponding to the PDF of the RVs:

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- What is a the structure of: $F \mathbf{v}=\mathbf{g}, F \in \mathbb{C}^{m w \times m w}, \mathbf{v}, \mathbf{g} \in \mathbb{C}^{m w}$

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- Not a "cheap" computation: direct matrix solvers, use an iterative method: $\mathrm{BiCgStab}+$ block-diagonal, mean-based pre-conditioner
- The matrix $F$ is never stored explicitly, instead we only need to be able to compute the product of $F$ with another vector for the iterative method to work
- Compute times and memory demands are still much higher than MC/SC. Scope for lot of research in sparse data structures for $F, \mathrm{~g}$

| SG-gPC | MC |
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- How do we choose the number of edges in the FEM formulation, $m$ ? Horizontally: Long enough to capture several correlation lengths Vertically: Sufficient for boundary conditions to be accurate
- How many basis functions to keep, w? Recall: $\Psi_{\mathbf{i}}(\vec{z})=\psi_{i_{1}}\left(z_{1}\right) \ldots \psi_{i_{d}}\left(z_{d}\right), \quad 0 \leq \mathbf{i}=\sum_{j=1}^{d} i_{j} \leq n$ (max degree) belong to the space $\mathbb{P}_{n}^{d}$ of dimension $w=\binom{n+d}{n}$ Choose $d$ : no of $K L$ terms till eigenvalue falls by $1 / 10$ Choose $n$ : depends on available computer resources



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Values of $d \mathrm{v} / \mathrm{s}$ surface length ( $a$ ) and correlation length ( $l$ )

| $\downarrow l \backslash \mathrm{a} \rightarrow$ | 15 | 30 |
| :---: | :---: | :---: |
| 1 | 15 | 29 |
| 0.5 | 29 | 58 |

## Compare Stochastic Galerkin with Monte Carlo

$$
h=0.05 /(2 \pi) \quad h=0.20 /(2 \pi)
$$




$$
\theta_{i}=40^{\circ}, \epsilon_{r}=4-j, a=20, l=1, \lambda=1 \quad d=15, n=1
$$

## Reflections

- SG-gPC is a deal breaker for even small values of $d$ $\rightarrow$ Computational innovation required
- For small-moderate values of $d$, (e.g. $d<50$ ) SC methods are a powerful alternative to MC $\rightarrow$ Can be improved by anisotropic sparse grids.
- For higher values of $d, \mathrm{MC}$ is the still optimal $\rightarrow$ Can be improved by Markov chain monte carlo


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## Thanks!

Most of this talk was based on:
"Stochastic Solutions to Rough Surface Scattering using the finite element method", Uday K Khankhoje and Shreyas Padhy; IEEE Transactions on Antennas and Propagation, 65(08) 2017.

Refer to http://www.ee.iitm.ac.in/uday for full-text of relevant papers. Or just Google my name.

E-mail: uday@ee.iitm.ac.in

## Aside: Why is Guass quadrature so effective?

- To approximate $\int f(x) w(x) d x$ by an $n$-point rule, express $f(x)$ in terms of: $n$th order polynomial $p_{n}(x)$, a quotient $q(x)$ (degree $<n$ ), \& a remainder $r(x)($ degree $<n)$ :

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\begin{equation*}
f(x)=q(x) * p_{n}(x)+r(x) \tag{1}
\end{equation*}
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- Integrate w.r.t. $w(x)$, and use orthogonality property of $p_{n}(x)$ : we get $\int f(x) w(x) d x=\int r(x) w(x) d x$. Next, construct a $n-1$ degree polynomial from $n$ points (Lagrange interpolation) and get:
- Here's where the choice of points $x_{i}$ is crucial: If $x_{i}$ is a root of $p_{n}(x)$, then from (1) it follows that $f\left(x_{i}\right)=r\left(x_{i}\right)$ and from (2) that
- Thus we say that Gauss quadrature is accurate to order $2 n-1$


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[^0]:    ${ }^{a}$ Khankhoje et al. "Computation of radar scattering from heterogeneous rough soil using the finite element method", 2013 IEEE TGRS
    ${ }^{b}$ Khankhoje et al. "Stochastic Solutions to Rough Surface Scattering using the finite element method", 2017 IEEE TAP

