

EE6340 - Information Theory

Problem Set 1 Solution

February 21, 2013

1. a) Random variable X = No. of coin tosses till the first head appears. If $\mathbb{P}(\text{head})=p$ and $\mathbb{P}(\text{tail}) = q$,

$$\begin{aligned} \mathbb{P}(x = n) &= pq^{n-1} \\ \implies H(X) &= - \sum_{n=1}^{\infty} pq^{n-1} \log(pq^{n-1}) \\ &= - \frac{p \log p}{1-q} - \frac{pq \log q}{p^2} \\ &= H(p)/p \end{aligned}$$

For $p = \frac{1}{2}$, $H(X) = H(0.5)/0.5 = 2$ bits.

- b) Best questions are those which have equal chances of being answered as Yes or as No.

$\mathbb{P}(x = 1) = \frac{1}{2}$ and $\mathbb{P}(x=2 \text{ or } 3 \text{ or } 4\dots) = \frac{1}{2}$ (equally probable).

So first question is, "Is it 1 or not?".

Similarly, $\mathbb{P}(x = 2) = \frac{1}{4}$ and $\mathbb{P}(x=3 \text{ or } 4 \text{ or } 5\dots) = \frac{1}{4}$ (equally probable)

So second question is, "Is it 2 or not?".

The questions proceed as above.

$\mathbb{E}(\text{No. of questions}) = \sum_{n=1}^{\infty} n \frac{1}{2^n} = 2 = H(X)$.

In general, $\mathbb{E}(\text{No. of questions}) \geq H(X)$

This problem can be interpreted as a source coding problem with 0=no, 1=yes, X =Source, Y =Encoded sequence.

2. a) $H(X, g(X)) = H(X) + H(g(X)|X)$ (Chain rule)
 b) $H(g(X)|X) = \sum_x p(x)H(g(X)|X = x) = \sum_x p(x)0 = 0$ (For a given x , $g(x)$ is fixed).
 $\implies H(X, g(X)) = H(X)$
 c) $H(X, g(X)) = H(g(X)) + H(X|g(X))$ (Chain rule)
 d) $H(X|g(X)) \geq 0$ with equality if $g(\cdot)$ is one-to-one.
 (a),(b) and (c) $\implies H(X, g(X)) \geq H(g(X))$
3. Let there be 2 y_i 's y_1 and y_2 such that for $x = x_0$, $p(x_0, y_1) > 0$ and $p(x_0, y_2) > 0$.
 $\implies p(y_1|x_0) > 0$ and $p(y_2|x_0) > 0$, neither of them being 0 or 1.

$$\begin{aligned} H(Y|X) &= - \sum_{x,y} p(x,y) \log_2 p(y|x) \\ &\geq -p(x_0)p(y_1|x_0) \log_2 p(y_1|x_0) - p(x_0)p(y_2|x_0) \log_2 p(y_2|x_0) \\ &> 0 \end{aligned}$$

since $-t \log t > 0$ for $0 < t < 1$. So, $H(Y|X) = 0$ iff Y is a function of X . Else $H(Y|X) > 0$.

4. X =Outcome of world series

Y =No.of games played $\in \{4,5,6,7\}$

For $Y = 4$, there are 2 outcomes $\{AAAA,BBBB\}$ each with probability $\frac{1}{2^4}$.

For $Y = 5$,there are $2 \times \binom{4}{3} = 8$ outcomes each with probability $\frac{1}{2^5}$.

For $Y = 6$,there are $2 \times \binom{5}{3} = 20$ outcomes each with probability $\frac{1}{2^6}$.

for $Y = 7$,there are $2 \times \binom{6}{3} = 40$ outcomes each with probability $\frac{1}{2^7}$.

Thus,

$$\begin{aligned}\mathbb{P}(Y = 4) &= \frac{1}{8} \\ \mathbb{P}(Y = 5) &= \frac{1}{4} \\ \mathbb{P}(Y = 6) &= \frac{5}{16} \\ \mathbb{P}(Y = 7) &= \frac{5}{16}\end{aligned}$$

X ={AAAA,BBBB,....,BABABAB,....BBBAAAA}

$\mathbb{P}(BAAAA) = \frac{1}{2^5}$, $\mathbb{P}(BABABAB) = \frac{1}{2^7}, \dots$

There are 2 sequences of length 4, 8 sequences of length 5, 20 sequences of length 6 and 40 sequences of length 7.

$$\begin{aligned}H(X) &= - \sum_x p(x) \log_2 p(x) \\ &= 2 \left(\frac{1}{16} \right) \log_2 16 + 8 \left(\frac{1}{32} \right) \log_2 32 + 20 \left(\frac{1}{64} \right) \log_2 64 + 40 \left(\frac{1}{128} \right) \log_2 128 \\ &= 5.8125 \\ H(Y) &= - \sum_y p(y) \log_2 p(y) \\ &= \frac{1}{8} \log_2 8 + \frac{1}{4} \log_2 4 + \frac{5}{16} \log_2 \frac{16}{5} + \frac{5}{16} \log_2 \frac{16}{5} \\ &= 1.924\end{aligned}$$

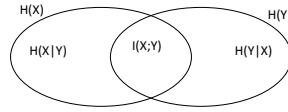
Y =length(X), i.e, Y is a deterministic function of X .

$$\begin{aligned}H(Y|X) &= 0 \\ H(X) + H(Y|X) &= H(X, Y) = H(Y) + H(X|Y) \\ \implies H(X|Y) &= H(X) - H(Y) \\ &= 3.889\end{aligned}$$

5. a) $H(X) = \frac{2}{3} \log_2 \frac{3}{2} + \frac{1}{3} \log_2 3 = 0.918 = H(Y)$
- b) $H(X|Y) = \frac{1}{3} H(X|Y=0) + \frac{2}{3} H(X|Y=1) = 0.667 = H(Y|X)$
- c) $H(X, Y) = H(X) + H(Y|X) = 1.585$
- d) $H(Y) - H(Y|X) = 0.251$
- e) $I(X; Y) = H(Y) - H(Y|X) = 0.251$
- f)

6. Identically distributed $\implies H(X_1) = H(X_2)$

$$\text{a) } \rho = 1 - \frac{H(X_2|X_1)}{H(X_1)} = \frac{H(X_1) - H(X_2|X_1)}{H(X_1)} = \frac{H(X_2) - H(X_2|X_1)}{H(X_1)} = \frac{I(X_1; X_2)}{H(X_1)}$$



- b) $I(X_1; X_2) = H(X_1) - H(X_1|X_2)$
 But $H(X_1|X_2) \geq 0 \implies I(X_1; X_2) \leq H(X_1) \implies \rho \leq 1$
 $I(X_1; X_2) \geq 0 \implies \rho \geq 0$
 $\therefore 0 \leq \rho \leq 1$
- c) X_1, X_2 are independent. So $H(X_1|X_2) = H(X_1) \implies I(X_1; X_2) = 0$
 Thus $\rho = 0$ when X_1 and X_2 are i.i.d
- d) $H(X_2|X_1) = 0$ when X_2 is a function of $X_1 \implies \rho = 1$.