

# EE5160: Error Control Coding

## Problem Set 7

1. Find the power series  $s(D)$  for the rational function

$$\frac{a(D)}{b(D)} = \frac{1 + D + D^2}{1 + D + D^3}$$

and note that, after a transient, its coefficients are periodic. Show that  $a(D) = b(D)s(D)$ .

2. Find  $G_{sys}(D)$ ,  $H_{sys}(D)$ ,  $G_{poly}(D)$  and  $H_{poly}(D)$  for the rate-2/3 convolutional-code generator matrix given by

$$G(D) = \begin{bmatrix} \frac{1+D}{1+D+D^2} & 1+D & 0 \\ 1 & \frac{1}{1+D} & \frac{D}{1+D^2} \end{bmatrix}.$$

3. Find the Type I realization of the transfer function

$$g_{ij}(D) = \frac{a_0 + a_1D + \dots + a_mD^m}{b_0 + b_1D + \dots + b_mD^m}.$$

To do so write  $c^{(j)}(D)$  as

$$c^{(j)}(D) = \left( u^{(i)}(D) \cdot \frac{1}{b(D)} \right) \cdot a(D).$$

Then sketch the direct implementation of the leftmost “filtering” operation  $v(D) = u^{(i)}(D) \cdot 1/b(D)$ , which can be determined from the difference equation

$$v_t = u_t^{(i)} - b_1v_{t-1} - b_2v_{t-2} - \dots - b_mv_{t-m}.$$

Next, sketch the direct implementation of the second filtering operation  $c^{(j)}(D) = v(D)a(D)$ , which can be determined from the difference equation

$$c_t^{(j)} = a_0v_t^{(i)} + a_1v_{t-1}^{(i)} \dots + a_mv_{t-m}^{(i)}.$$

Finally, sketch the two “filters” in cascade, the  $1/b(D)$  filter followed by the  $a(D)$  filter (going from left to right), and notice that the  $m$  delay (memory) elements may be shared by the two filters.

4. Consider the rate-2/3 convolutional code with

$$G(D) = \begin{bmatrix} \frac{1+D}{1+D+D^2} & 0 & 1+D \\ 1 & \frac{1}{1+D} & \frac{D}{1+D^2} \end{bmatrix}.$$

Let the input for the encoder  $G(D)$  be  $u(D) = [1 \quad 1+D]$  and find the corresponding codeword  $c(D) = [c_1(D) \quad c_2(D) \quad c_3(D)]$ . Find the input which yields the same codeword when the encoder is given by  $G_{sys}(D)$ . Repeat for  $G_{poly}(D)$ .

5. Consider a rate-2/3 convolutional code with

$$G_{poly}(D) = \begin{bmatrix} 1+D & 0 & 1 \\ 1+D^2 & 1+D & 1+D+D^2 \end{bmatrix}.$$

Show that the memory required for the Type I and TypeII realizations of  $G_{poly}(D)$  is  $\mu = 3$  and  $\mu = 6$ , respectively. Show that

$$G_{sys}(D) = \begin{bmatrix} 1 & 0 & \frac{1}{1+D} \\ 0 & 1 & \frac{D^2}{1+D} \end{bmatrix},$$

and that the memory required for its Type I realization is  $\mu = 3$ . Finally, show that the Type II realization of  $G_{sys}(D)$  requires only  $\mu = 2$ . Thus, the Type II realization of  $G_{sys}(D)$  is the minimal encoder for rate  $k/k + 1$  convolutional codes.

6. Show that

$$\max(x, y) = \log \left( \frac{e^x + e^y}{1 + e^{-|x-y|}} \right).$$

Hint: First suppose  $x > y$ .

7. Draw the state diagram for the rate-1/2 encoder described by

$$u(D) = [ 1 + D + D^2 \quad 1 + D^2 ].$$

Find the input output weight enumerator (IO-WE)

$$A'(I, W) = \sum_{i,w} A'_{i,w} I^i W^w$$

where  $A'_{i,w}$  is the number of weight- $w$  paths, corresponding to weight- $i$  encoder inputs, that diverge from the all-zeros trellis path one time in  $L$  trellis stages.

8. **(Optional)** Assuming the BI-AWGN, simulate Viterbi decoding of the rate-1/2 convolutional code whose encoder matrix is given by

$$G(D) = [ 1 + D^2 + D^3 + D^4 \quad 1 + D + D^4 ].$$

Plot the bit error rate from  $P_b = 10^{-1}$  to  $P_b = 10^{-6}$ . Repeat for the BCJR decoder and comment on the performance of the two decoders.