

Antenna diversity: (or spatial diversity).

— Obtained by placing multiple antennas at the transmitter and/or the receiver.

— Channels between different pairs of antennas are independent if

* antennas are spaced sufficiently apart
This ^{spacing} depends on

- scattering environment
- carrier frequency

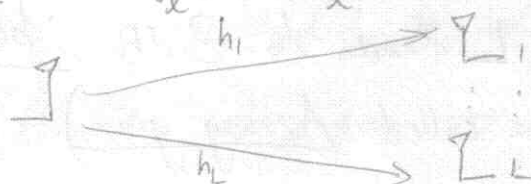
— More about MIMO channel modelling in Chapter 7.
For now, we will assume that independent channels can be achieved.

We will discuss the following:

- (1) Receive diversity [Single Input Multiple Output channel]
(SIMO)
- (2) Transmit diversity [Multiple Input Single Output channel]
(MISO)
- (3) 2x2 MIMO [2x2 Multiple Input Multiple Output channel]
(MIMO)
Two Input. Two Output

Receive diversity: (common in uplink)

Model: $y_l[m] = h_l[m] x[m] + w_l[m], l=1, \dots, L$



$w_e[m] \sim \text{CN}(0, N_0)$ independent across antennas.

$h_e[m] \sim$ independent Rayleigh fading

Diversity gain = L . (Similar to Time Diversity analysis)

$x[i]$ is detected based on $y_1[i], y_2[i], \dots, y_L[i]$.

$\hat{x}[i]$ can be detected from $\underline{h}^H \underline{y}$ as before.

BPSK : $(\pm a)$

$$P_e(\text{error} | \underline{h}) = Q\left(\sqrt{2\|\underline{h}\|^2 \text{SNR}}\right) \quad \text{where } \text{SNR} = \frac{a^2}{N_0}$$

Received SNR conditioned on the channel \underline{h} is

$$\|\underline{h}\|^2 \text{SNR}.$$

This can be thought of as consisting of two parts

$$\left(\frac{1}{L} \|\underline{h}\|^2\right) (L \text{SNR})$$

* By adding multiple receive antennas, received power increases (linearly with L). This is reflected in $L \text{SNR}$.

(In the time-diversity repetition coding, received power increases with more transmissions only because transmit power increases linearly. If total transmit power per symbol is kept constant, receive power will not increase.)

\Rightarrow Doubling L leads to 3 dB power gain.

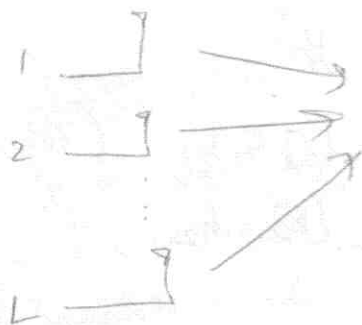
(Also called Array gain)

* $\frac{1}{L} \| \underline{h} \|^2$ represents the averaging over multiple independent paths.

$$\frac{1}{L} \| \underline{h} \|^2 = \frac{1}{L} \sum_{l=1}^L |h_l|^2$$

As L increases, due to law of large numbers, $\frac{1}{L} \| \underline{h} \|^2$ converges to 1 and P_r (deep fade) reduces. This gives us the exponent L in the P_r (error) or the Diversity gain.

Transmit diversity: (common in downlink)



Model:

$$y[n] = \sum_{l=1}^L h_l[n] x_l[n] + w[n]$$

\uparrow
 $\mathcal{N}(0,1)$

* Simple method to get a diversity gain of L .

- Transmit the same symbol over the L different antennas during L symbol times. At any time, only one antenna is used. (Repetition code, does not use all available degrees of freedom).

- Any time diversity code of block length L can be used: simply use one antenna at a time and transmit the coded symbols ^{successively over the diff. antennas} one antenna at a time (Provides coding gain over the repetition code).

* One can also design codes specifically for the transmit diversity system. These codes are referred to as "space-time codes". Codewords span across antennas and time.

* Simple space-time codes. (More MIMO in chaps. 7-10).

Alamouti scheme:

Consider $L=2$.

$$y[m] = h_1[m] x_1[m] + h_2[m] x_2[m] + w[m].$$

* Transmits symbols u_1 and u_2 over 2 symbol times.

$$\begin{aligned} x_1[1] &= u_1 & x_1[2] &= -u_2^* \\ x_2[1] &= u_2 & x_2[2] &= u_1^* \end{aligned}$$

Lecture 15 (3/9/2008)

* If we assume $h_1[1] = h_1[2] = h_1$ (coh. time \gg 2 symbols)
 $h_2[1] = h_2[2] = h_2$,

$$\begin{bmatrix} y[1] & y[2] \end{bmatrix} = \begin{bmatrix} h_1 & h_2 \end{bmatrix} \begin{bmatrix} u_1 & -u_2^* \\ u_2 & u_1^* \end{bmatrix} + \begin{bmatrix} w[1] & w[2] \end{bmatrix}$$

Re-write as:

$$\begin{bmatrix} y[1] \\ y^*[2] \end{bmatrix} = \begin{bmatrix} h_1 & h_2 \\ h_2^* & -h_1^* \end{bmatrix} \begin{bmatrix} u_1 \\ u_2 \end{bmatrix} + \begin{bmatrix} w[1] \\ w^*[2] \end{bmatrix}$$

\uparrow
 H : Columns are orthogonal

\Rightarrow detection of u_1, u_2 decomposes into 2 separate scalar problems.

$$f(\underline{y}|\underline{u}, b) = \text{CN}(\underline{H}\underline{u}, N_0 \underline{I})$$

(26)

$$\|\underline{y} - \underline{H}\underline{u}\|^2 = (\underline{y} - \underline{H}\underline{u})^H (\underline{y} - \underline{H}\underline{u})$$

$$= \underline{y}^H \underline{y} - \underline{y}^H \underline{H}\underline{u} - \underline{u}^H \underline{H}^H \underline{y} + \underline{u}^H \underline{H}^H \underline{H}\underline{u}$$

$\underline{H}^H \underline{y}$ is sufficient for detection.

$$\underline{H}^H \underline{y} = \underline{H}^H \underline{H}\underline{u} + \underline{H}^H \underline{w}$$

$$(H \text{ is orthogonal}) (H^H H = (|h_1|^2 + |h_2|^2) \underline{I})$$

$$\underline{r} = \underline{H}^H \underline{y} = \underline{u} (|h_1|^2 + |h_2|^2) + \underline{w}'$$

$$\underline{w}' \sim \text{CN}(\underline{0}, N_0 (|h_1|^2 + |h_2|^2) \underline{I})$$

$$r_1 = (|h_1|^2 + |h_2|^2) u_1 + w_1$$

$$= \| \underline{h} \|^2 u_1 + w_1 \leftarrow \text{Diversity 2}$$

$$r_2 = \| \underline{h} \|^2 u_2 + w_2 \leftarrow$$

* In the repetition code, ~~one~~ symbol is transmitted over 2 symbol intervals. Here, 2 symbols are transmitted over 2 symbol intervals.

* To maintain the same transmit power in both cases, each symbol is transmitted with half the power in each interval for the Alamouti scheme.

* Alamouti scheme works for any constellation.

Suppose, we use BPSK. \rightarrow 2 bits over 2 symbol times

For the same rate, repetition code should use 4-PAM.

To achieve same d_{\min} , 4-PAM needs 5 times the energy as BPSK

For the same transmit power, factor of 2 more power is used on each symbol in repetition code.

Overall, repetition requires 2.5 times more power than Alamouti scheme ($\approx 4 \text{ dB}$).

- * Repetition uses only one degree of freedom of the channel over the 2 symbol durations
→ along the dimension $[h_1, h_2]$.

Alamouti scheme uses 2 degrees of freedom available in the channel over the 2 symbol durations
→ along dimensions $[h_1, h_2^*]$

and $[h_2, -h_1^*]$.

Lecture 16: (9 Sep 2008)

Space-time code design:

- design a good code to
- * In order to exploit time diversity, we had to maximize the minimum product distance between codewords.
 - * We can show a similar criterion for space-time codes.

- A codeword in a space-time code is a matrix

X_i : $L \times N$ matrix

L antennas, N time intervals
(block length = N).

A space-time code is specified by a set of codeword matrices $\{X_i\}$.

Eg: Alamouti scheme

X_i is of the form $\begin{bmatrix} u_1 & -u_2^* \\ u_2 & u_1^* \end{bmatrix}$.

$L=2$
 $N=2$

Repetition scheme

X_i is of the form $\begin{bmatrix} u & 0 \\ 0 & u \end{bmatrix}$

- Assume that the channel remains the same for N time intervals (27)

$$\begin{bmatrix} y[1] & \dots & y[N] \end{bmatrix} = \begin{bmatrix} h_1 & h_2 & \dots & h_L \end{bmatrix} \begin{bmatrix} X \end{bmatrix} + \begin{bmatrix} w[1] & \dots & w[N] \end{bmatrix}$$

$$\underline{y}^T = \underline{h}^H X + \underline{w}^T$$

$$\underline{h} = \begin{bmatrix} h_1^* \\ h_2^* \\ \vdots \\ h_L^* \end{bmatrix}$$

Normalize the codewords such that the average energy per symbol time is 1. $\Rightarrow \text{SNR} = \frac{1}{N_0}$.

• Pairwise error probability $P_r(X_A \rightarrow X_B)$

$$P_r(X_A \rightarrow X_B | \underline{h}) \leq Q \left(\frac{\| \underline{h}^H (X_A - X_B) \|}{2\sqrt{N_0/2}} \right)$$

$$P_r(X_A \rightarrow X_B) \leq E \left[Q \left(\frac{\| \underline{h}^H (X_A - X_B) \|}{2\sqrt{N_0/2}} \right) \right]$$

$$= E \left[Q \left(\sqrt{\frac{\text{SNR} \underline{h}^H (X_A - X_B) (X_A - X_B)^H \underline{h}}{2}} \right) \right]$$

$(X_A - X_B)(X_A - X_B)^H$ is Hermitian (A complex square matrix X is Hermitian if $X = X^H$)

and can be diagonalized as

$U \Lambda U^H$, where U is unitary.

$\Lambda = \text{diag}(\lambda_1^2, \lambda_2^2, \dots, \lambda_L^2)$ where λ_i are the singular values of $X_A - X_B$.

$$P_r(X_A \rightarrow X_B) \leq E \left[Q \left(\sqrt{\frac{\text{SNR} \sum_{l=1}^L |\tilde{h}_l|^2 \lambda_l^2}{2}} \right) \right]$$

where $\tilde{\underline{h}} = \underline{U}^H \underline{h}$.

$\tilde{\underline{h}}$ has the same distribution as \underline{h} .

$|\tilde{h}_l|^2$ is exponential.

$$\Rightarrow P_r(X_A \rightarrow X_B) \leq \prod_{l=1}^L \frac{1}{1 + \frac{\text{SNR} \lambda_l^2}{4}}$$

— If all λ_l^2 are > 0 for all codeword differences $(X_A - X_B)$, maximal diversity gain of L is achieved. This is the rank criterion.

Num. of positive eigen values = rank of $(X_A - X_B) \leq N$

Therefore, full diversity is possible only for $N \geq L$.

— If all $\lambda_l^2 > 0$, then

$$P_r(X_A \rightarrow X_B) \leq \frac{4^L}{\text{SNR}^L \prod_{l=1}^L \lambda_l^2} \quad (\text{at high SNR})$$

↓
determinant of $(X_A - X_B)(X_A - X_B)^H$

Therefore, the coding gain is determined by the minimum of the determinant $\det[(X_A - X_B)(X_A - X_B)^H]$ over all codeword pairs X_A, X_B .

Therefore, we should maximize the minimum determinant. This is the determinant criterion.