

# LECTURE-8

Note Title

28-Jun-19

## CONVEX - CONCAVE    MINIMAX    OPTIMIZATION

$$\min_{x \in X} \max_{y \in Y} f(x, y)$$

$f(\cdot, y)$  is convex

$f(x, \cdot)$  is concave

$$\left\{ \begin{array}{l} g \text{ is concave} \\ \triangleq -g \text{ is convex} \\ 1) g(\alpha y_1 + (1-\alpha)y_2) \geq \alpha g(y_1) + (1-\alpha)g(y_2) \\ 2) g(z) \leq g(y) + \langle \nabla g(y), z - y \rangle \end{array} \right.$$

i) Constrained Optimization :

$$\min_{\boldsymbol{x}} \max_{i=1, \dots, m} f_i(\boldsymbol{x}) \quad \leftarrow$$

$$\begin{aligned} & \text{Find } \boldsymbol{x} \\ & \text{s.t. } f_i(\boldsymbol{x}) \leq 0 \quad i=1, \dots, m. \end{aligned}$$

Each  $f_i$  is convex.

III

$$\begin{aligned} & \min_{\boldsymbol{x}} \max_{\substack{\boldsymbol{y} \\ \boldsymbol{y}_i \geq 0 \\ \sum_{i=1}^m \boldsymbol{y}_i = 1}} \sum_{i=1}^m \boldsymbol{y}_i f_i(\boldsymbol{x}) \end{aligned}$$

$$f(\boldsymbol{x}, \boldsymbol{y}) = \sum_{i=1}^m \boldsymbol{y}_i f_i(\boldsymbol{x})$$

$f(\cdot, \boldsymbol{y})$  is convex if  $\boldsymbol{y}$   
 $f(\boldsymbol{x}, \cdot)$  is linear in  $\cdot$  if  $\boldsymbol{x}$  concave

2) Many non-smooth functions can be written as smooth minimax problems.

$$a) \quad \|x\|_1 = \max_{-1 \leq y_i \leq 1} \sum_i y_i x_i$$

$$\min_x \frac{1}{2} \|Ax - b\|^2 + d \|x\|_1 = \min_x \max_{-1 \leq y_i \leq 1} \underbrace{\frac{1}{2} \|Ax - b\|^2 + d \sum_i y_i x_i}_{\text{---}}$$

b)

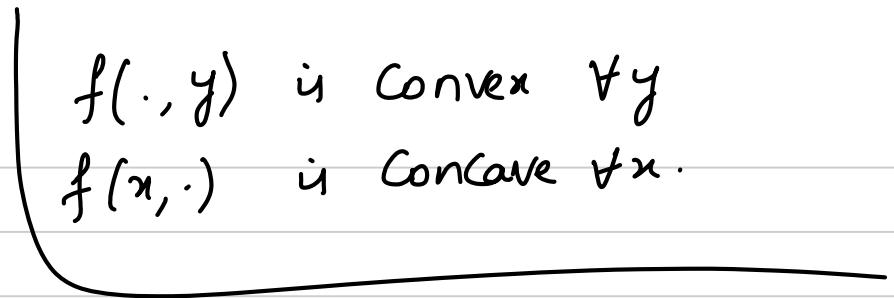
Given  $a_1, \dots, a_m$  find the  
smallest ball covering  $a_1, \dots, a_m$ .

→ Again can be written as smooth  
minimax problem.

3) Minimax ≡ zero sum games.

$$\min_x \underbrace{\max_y f(x, y)}_{\equiv g(x)}$$

$f(\cdot, y)$  is convex w.r.t  $y$   
 $f(x, \cdot)$  is concave w.r.t  $x$ .



Exercise:  $g(x)$  is convex

How to evaluate  $\nabla g(x)$ ?

(Informal) Danskin's thm: Under smoothness assumptions on  $f(\cdot, \cdot)$ ,

$$\nabla g(x) = \text{Conv. hull} \left\{ \nabla_x f(x, y) : y \in \arg \underline{\max}_z f(x, z) \right\}.$$

$$\left. \begin{array}{l} 1) \text{ Find } y_t \in \arg \max_y f(x_t, y) \\ 2) \quad x_{t+1} = x_t - \eta \nabla_x f(x_t, y_t) \end{array} \right\} \begin{array}{l} \text{Convex opt.} \\ \text{Subgradient descent} \\ \text{on } g(x). \end{array}$$

Issues : ① Finding  $y_t$  takes time  
 ②  $g$  could be non smooth  $\Rightarrow$  slow convergence rate.

Non smooth : Gradient descent ascent  $\rightarrow O(\frac{1}{\sqrt{T}})$

Smooth : Mirror-Prox  $\rightarrow O(\frac{1}{T})$ .

$g(x)$  could still be non smooth

[ $f$  could be arbitrary] maximin  $\leq$  minimax

Weak duality:  $\underbrace{\max_{y \in Y} \min_{x \in X} f(x, y)}_{LHS} \leq \underbrace{\min_{x \in X} \max_{y \in Y} f(x, y)}_{RHS}.$

$$y^* = \arg \max_{y \in Y} \left[ \min_{x \in X} f(x, y) \right]$$

$$x^* = \arg \min_{x \in X} \left[ \max_{y \in Y} f(x, y) \right]$$

$$LHS = \min_{x \in X} f(x, y^*) \leq f(x^*, y^*) \leq \max_{y \in Y} f(x^*, y) = RHS.$$

[Sion's minimax thm.]

Strong duality: If  $f(x, y)$  is Convex-Concave and  $X$  and  $Y$

are Compact then  $\min_{x \in X} \max_{y \in Y} f(x, y) = \max_{y \in Y} \min_{x \in X} f(x, y)$ .

---

Defn.  $\epsilon$ -primal dual pair:  $(\bar{x}, \bar{y})$  is Said to be an  $\epsilon$ -primal

dual pair if  $\max_y f(\bar{x}, y) - \min_x f(x, \bar{y}) \leq \epsilon$ .

Note:  $\max_y f(\bar{x}, y) \geq f(\bar{x}, \bar{y}) \geq \min_x f(x, \bar{y})$ .

Let  $(\bar{x}, \bar{y})$  be an  $\epsilon$ -primal dual pair

Exercise: Show that  $\bar{x}$  is an  $\epsilon$ -optimal soln. for

$$\min_x g(x) \triangleq \max_y f(x, y)$$

Exercise:  $\bar{y}$  is an  $\epsilon$ -optimal soln. for  $\max_y h(y) \triangleq \min_x f(x, y)$ .

$$\text{diam}(X), \text{diam}(Y) \leq R$$

Gradient descent ascent :

$$x_{t+1} = x_t - \eta \cdot \nabla_x f(x_t, y_t) \rightarrow \text{GD for min}$$

$$y_{t+1} = y_t + \eta \cdot \nabla_y f(x_t, y_t) \rightarrow \text{GA for max}$$

$f(\cdot, y)$  is convex

$f(x, \cdot)$  is concave

$$\|\nabla_x f(x, y)\| \leq G$$

$$\|\nabla_y f(x, y)\| \leq G$$

Thm: GDA has Conv. rate  $O\left(\frac{\epsilon R}{\sqrt{T}}\right)$ .

Proof:  $\|x_{t+1} - x\|^2 = \|x_t - x\|^2 - 2\eta \underbrace{\langle \nabla_x f(x_t, y_t), x_t - x \rangle}_{\text{Conv. of } f(\cdot, y_t)} + \eta^2 \underbrace{\|\nabla_x f(x_t, y_t)\|^2}_{G\text{-Lipschitz}}$

$x$  is arb.

$$\leq \|x_t - x\|^2 - 2\eta [f(x_t, y_t) - f(x, y_t)] + \eta^2 G^2$$

$$\boxed{y \text{ in arb.}} \quad \|y_{t+1} - y\|^2 = \|y_t - y\|^2 - 2\eta \underbrace{\langle \nabla_y f(x_t, y_t), y - y_t \rangle}_{\text{Concavity of } f(x_t, \cdot)} + \eta^2 \underbrace{\|\nabla_y f(x_t, y_t)\|^2}_{L\text{-Lipschitz}}$$

$$\leq \|y_t - y\|^2 - 2\eta [f(x_t, y) - f(x_t, y_t)] + \eta^2 \varsigma^2$$

$$\|x_{t+1} - x\|^2 + \|y_{t+1} - y\|^2 \leq \|x_t - x\|^2 + \|y_t - y\|^2 - 2\eta [f(x_t, y) - f(x, y_t)] + 2\eta^2 \varsigma^2.$$

$$\frac{1}{T+1} \sum_{t=0}^T [f(x_t, y) - f(x, y_t)] \leq \frac{\|x_0 - x\|^2 + \|y_0 - y\|^2}{2\eta(T+1)} + \eta \varsigma^2$$

$$f\left(\underbrace{\frac{1}{T+1} \sum_{t=0}^T x_t}_{\bar{x}_T}, y\right) \leq \frac{1}{T+1} \sum_{t=0}^T f(x_t, y) \quad \text{and} \quad f\left(x, \underbrace{\frac{1}{T+1} \sum_{t=0}^T y_t}_{\bar{y}_T}\right) \geq \frac{1}{T+1} \sum_{t=0}^T f(x, y_t)$$

$$f(\bar{x}_T, y) - f(x, \bar{y}_T) \leq \frac{\|x_0 - x\|^2 + \|y_0 - y\|^2}{2\eta(T+1)} + \gamma \zeta^2.$$

$$\begin{aligned} \max_y f(\bar{x}_T, y) - \min_x f(x, \bar{y}_T) &\leq \frac{\max_x \|x_0 - x\|^2 + \max_y \|y_0 - y\|^2}{2\eta(T+1)} + \gamma \zeta^2 \\ &\leq \frac{R^2}{\eta(T+1)} + \gamma \zeta^2 \\ \gamma = \frac{R}{6\sqrt{T+1}} &\leq \frac{2GR}{\sqrt{T+1}}. \end{aligned}$$

$(\bar{x}_T, \bar{y}_T)$  is an  $\frac{2GR}{\sqrt{T+1}}$  - primal dual pair.

⑩

[Prox] Proximal algorithm:  $x_{t+1} = \arg \min_{x \in X} \left[ f(x) + \frac{1}{2\eta} \|x - x_t\|^2 \right]$

GD Step:  $x_{t+1} = \arg \min_{x \in X} \left[ f(x_t) + \langle \nabla f(x_t), x - x_t \rangle + \frac{1}{2\eta} \|x - x_t\|^2 \right]$

Exercise: Prox has convergence rate  $O(\frac{1}{T})$  for  
[possibly non-smooth] convex functions.

Exercise: Prox step is efficiently implementable for

$L$ -Smooth functions and  $\eta \leq \frac{1}{2L}$ .

Efficient: Can find  $\epsilon$ -approximate Soln. to the prox step  
in  $O(\log \frac{1}{\epsilon})$  iterations.

Prox Step: Given  $x$ , find  $w^*$  s.t.

$$w^* = \underset{z \in X}{\operatorname{arg\,min}} \quad f(z) + \frac{1}{2\eta} \|x - z\|^2$$

$$\equiv \quad w^* = \quad x - \eta \nabla f(w^*) \rightarrow \text{Prox.}$$

$$\nabla f(x) \rightarrow \text{GD}$$

Algorithm to find ]. Pick  $w_0 = x$   
 From Step  $w_{t+1} = x - \eta \nabla f(w_t) \rightarrow ① //$

Ⓐ Claim:  $w^*$  is the unique fixed point of ①

Ⓑ Claim: ① is a  $\frac{1}{2}$  Contraction.

$$\begin{aligned} w \rightarrow w^+ &= x - \eta \nabla f(w) \\ \tilde{w} \rightarrow \tilde{w}^+ &= x - \eta \nabla f(\tilde{w}) \end{aligned} \quad \left. \begin{aligned} \|w^+ - \tilde{w}^+\| &= \eta \|\nabla f(w) - \nabla f(\tilde{w})\| \\ &\leq \eta L \|w - \tilde{w}\| \end{aligned} \right\}$$

If  $\eta \leq \frac{1}{2L}$   $\leq \frac{1}{2} \|w - \tilde{w}\|$ .

Ⓐ + Ⓑ  $\Rightarrow$  ① finds  $\epsilon$ -approximate  $w^*$  in  $\log \frac{1}{\epsilon}$  iterations.

## Conceptual Mirror-Prox

$$x_{t+1} = \underset{x \in X}{\operatorname{argmin}} f(x, y_{t+1}) + \frac{1}{2\eta} \|x - x_t\|^2$$

$$y_{t+1} = \underset{y \in Y}{\operatorname{argmax}} f(x_{t+1}, y) - \frac{1}{2\eta} \|y - y_t\|^2$$

Implementation of each step

$$\begin{aligned} x_t^0 &= x_t \\ y_t^0 &= y_t \end{aligned}$$

$$i=1, \dots, \log \frac{1}{\epsilon}, \quad x_t^i = \underset{x \in X}{\operatorname{argmin}} f(x, y_t^{i-1}) + \frac{1}{2\eta} \|x - x_t\|^2$$

*2' in the actual mirror prox alg.*

$$y_t^i = \underset{y \in Y}{\operatorname{argmax}} f(x_t^{i-1}, y) - \frac{1}{2\eta} \|y - y_t\|^2.$$

Exercise: For Conceptual Mirror-Prox, obtain  $O\left(\frac{1}{\tau}\right)$  convergence rate for L-Smooth minimax opt.

Exercise: Show that the algorithm for implementing the Prox Step Converges in  $O\left(\log \frac{1}{\epsilon}\right)$  iterations.

Exercise: Careful accounting of all the terms to show that 2 inner steps suffice for  $O\left(\frac{1}{\tau}\right)$  convergence rate.