

Assignment 1

Basics, Source coding, RLL codes

Date Assigned: Feb 9

1. *Example of Joint entropy.* Let $p(x, y)$ given by

$X \setminus Y$	0	1
0	1/3	1/3
1	0	1/3

Find

- (a) $H(X), H(Y)$
 - (b) $H(X|Y), H(Y|X)$
 - (c) $H(X, Y), I(X; Y)$
2. Let $X, Y,$ and Z be discrete random variables with a joint distribution. Prove the validity of the following inequalities and if true find conditions for equality:
- (a) $I(X, Y; Z) \geq I(X; Z).$
 - (b) $H(X, Y|Z) \geq H(X|Z).$
 - (c) $H(X, Y; Z) \geq I(X; Z).$
 - (d) $I(X; Z|Y) \geq I(Z; Y|X) - I(Z; Y) + I(X; Z).$
 - (e) $H(X, Y, Z) - H(X, Y) \leq H(X, Z) - H(X).$

3. *A measure of correlation.* Let X_1 and X_2 be identically distributed, but not necessarily independent. Let

$$\rho = 1 - \frac{H(X_2|X_1)}{H(X_1)}$$

- (a) Show $\rho = I(X_1; X_2)/H(X_1).$
 - (b) Show $0 \leq \rho \leq 1$
 - (c) When is $\rho = 0$? When is $\rho = 1$?
4. Let X, Y and Z be binary random variables so that the eight elements in the joint $\{X, Y, Z\}$ ensemble can be taken as the vertices of a unit cube.
- (a) Find a joint probability assignment $p(x, y, z)$ such that $I(X; Y) = 0$ and $I(X; Y|Z) = 1$ bit.
 - (b) Find the joint probability assignment $p(x, y, z)$ such that $I(X; Y) = 1$ bit and $I(X; Y|Z) = 0.$

The point of the problem is that no general inequality exists between $I(X; Y)$ and $I(X; Y|Z).$

5. *Conditional Mutual Information.* Consider a sequence of n binary random variables $X_1, X_2, X_3, \dots, X_n.$ Each sequence with an even number of 1's has the probability $2^{-(n-1)}$ and each sequence with odd number of 1's has the probability 0. Find the mutual informations $I(X_1; X_2), I(X_2; X_3|X_1), \dots, I(X_{n-1}; X_n|X_1, \dots, X_{n-2}).$
6. *The AEP and source coding.* A discrete memoryless source emits a sequence of statistically independent binary digits with probabilities $p(1) = 0.005$ and $p(0) = 0.995.$ The digits are taken 100 at a time and a binary codeword is provided for every sequences of 100 digits containing three or fewer ones.
- (a) Assuming that all codewords are the same length, find the minimum length required to provide codewords for all sequence with three or fewer ones.

- (b) Calculate the probability of observing a source sequence for which no codeword has been assigned.

7. *Entropy rates of Markov chains.*

- (a) Find the entropy rate of the two-state Markov chain with transition matrix

$$P = \begin{bmatrix} 1 - p_{01} & p_{01} \\ p_{10} & 1 - p_{10} \end{bmatrix}$$

- (b) What values of p_{01} and p_{10} maximize the rate of Part (a)?
 (c) Find the entropy rate of the two-state Markov chain with transition matrix

$$P = \begin{bmatrix} 1 - p & p \\ 1 & 0 \end{bmatrix}$$

- (d) Find the maximum value of the entropy rate of the Markov chain in Part (c). We expect that the maximizing value of p should be less than $1/2$, since the 0 state permits more information to be generated than the 1 state.

8. *Initial conditions.* For a Markov chain $\{X_n\}$, show that

$$H(X_0|X_n) \geq H(X_0|X_{n-1})$$

Thus initial conditions X_0 is become more difficult to recover as the future X_n unfolds.

9. *Pairwise independence.* Let X_1, X_2, \dots, X_{n-1} be *i.i.d* random variables taking values in $\{0,1\}$, with $\Pr\{X_i = 1\} = 1/2$. Let $X_n = 1$ if $\sum_{i=1}^{n-1} X_i$ is odd and $X_n = 0$ otherwise. Let $n \geq 3$.

- (a) Show that X_i and X_j are independent for $i \neq j, i, j \in \{1, 2, \dots, n\}$.
 (b) Find $H(X_i, X_j)$ for $i \neq j$.
 (c) Find $H(X_1, X_2, \dots, X_n)$. Is this equal to $nH(X_1)$?

10. *Huffman coding.* Consider the random variable

$$X = \begin{pmatrix} x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 \\ 0.49 & 0.26 & 0.12 & .04 & 0.04 & 0.03 & 0.02 \end{pmatrix}$$

- (a) Find a binary Huffman code for X.
 (b) Find the expected codelength for this encoding.
 (c) Find a ternary Huffman code for X.

11. *More Huffman codes.* Find the binary Huffman code for the source with probabilities $(1/3, 1/5, 1/5, 2/15, 2/15)$. Argue that this code is also optimal for the source with probabilities $(1/5, 1/5, 1/5, 1/5, 1/5)$.

12. *Bad codes.* Which of these codes cannot be Huffman code for any probability assignment?

- (a) $\{0,10,11\}$
 (b) $\{00,01,10,110\}$
 (c) $\{01,10\}$

13. Consider a random variable X which takes on four values with probabilities $(1/3, 1/3, 1/4, 1/12)$.

- (a) Construct a Huffman code for this random variable
 (b) Show that there exist two different sets of optimal length for the codewords, namely, show that codeword length assignments $(1,2,3,3)$ and $(2,2,2,2)$ are both optimal.

14. *Huffman code.* Find the (a) binary and (b) ternary Huffman codes for the random variable X with probabilities

$$p = \left(\frac{1}{21}, \frac{2}{21}, \frac{3}{21}, \frac{4}{21}, \frac{5}{21}, \frac{6}{21} \right).$$

- (c) Calculate $L = \sum p_i l_i$ in each case.

15. *Classes of codes.* Consider the code $\{0,01\}$

- (a) Is it instantaneous?
- (b) Is it uniquely decodable?
- (c) Is it nonsingular?

16. A source has an alphabet of 4 letters. The probabilities of the letters and two possible sets of binary codeword for the source are given below:

Letter	Prob.	Code I	Code II
a_1	0.4	1	1
a_2	0.3	01	10
a_3	0.2	001	100
a_4	0.1	000	1000

For each code, answer the following questions (no proofs or numerical answers are required).

- (a) Does the code satisfy the prefix condition?
 - (b) Is the code uniquely decodable?
 - (c) What is the mutual information provided about the event that the source letter is a_1 by the event that the first letter of the codeword is 1?
 - (d) What is the average mutual information provided about the source letter by the specification of the first letter of the codeword? Give a heuristic description of the purpose of the first letter in the codewords of Code II.
17. A discrete memoryless source with alphabet $A = \{a, b, c\}$ emits the following string.

b c c a c b c c c c c c c c c c a c c c a

Using the Lempel-Ziv algorithm, encode this sequence and find the code dictionary and transmitted sequence.

18. A source with alphabet $A = \{a, b, c\}$ is encoded using the Lempel-Ziv algorithm. The transmitted codeword sequence is

2, 3, 3, 1, 3, 4, 5, 10, 11, 6, 10.

Construct the dictionary and decode the sequence.