

Non binary BCH codes

Note Title

Ex: $n=10$, $q=9$

$$GF(9) = \{a + b\alpha : a, b \in GF(3)\}$$

$\beta \in GF(81)$, primitive

$$\alpha^2 + \alpha + 2 = 0$$

$$n=10 \mid \text{ord}(\beta) = 80$$

$$\{0\}, \{5\},$$

$$\Rightarrow \text{ord}(\beta^8) = n = 10 \quad \{1, 9\}, \{2, 8\}, \{3, 7\},$$

$$x^{10} - 1 = (x - 1)(x + 1)(x - \beta^8)(x - \beta^{72}) \dots$$

Ex: $q=4$, $n=15$

$$\alpha \in GF(16), \text{primitive}, \alpha^4 + \alpha + 1 = 0$$

$$\{0\}, \{1, 4\}, \{2, 8\}, \{3, 12\}, \{5\}, \{6, 9\}, \{10\}$$

$$\{7, 13\}, \{11, 14\}$$

$$M_\alpha(x) = x^2 + x + \alpha^5$$

$$x^{15} + 1 = (x + 1)(x + \alpha^5)(x + \alpha^{10}) \underbrace{(x + \alpha)(x + \alpha^4)}_{M_\alpha(x)}$$

$$M_{\alpha^2}(x) = \underbrace{(x + \alpha^2)(x + \alpha^8)}_{x^2 + x + \alpha^{10}} \underbrace{(x + \alpha^3)(x + \alpha^{12})}_{x^2 + \alpha^{10}x + 1} \underbrace{(x + \alpha^6)(x + \alpha^9)}_{x^2 + \alpha^5x + 1}$$

$$\frac{(x + \alpha^7)(x + \alpha^{13})}{x^2 + \alpha^5 x + \alpha^5} \quad \frac{(x + \alpha^{11})(x + \alpha^{19})}{x^2 + \alpha^{10} x + \alpha^{10}}$$

Possible $g(x)$: $(x^2 + x + \alpha^5)(x^2 + x + \alpha^{10})$

BCH codes over $GF(4)$ with $n=15$

$$t=1 \quad g(x) = \text{LCM}(M_\alpha(x), M_{\alpha^2}(x)) = x^4 + x + 1, k=11$$

$$t=2 \quad g(x) = \text{LCM}(M_\alpha(x), M_{\alpha^2}(x), M_{\alpha^3}(x), M_{\alpha^4}(x))$$

$$= (x^2 + x + \alpha^5)(x^2 + x + \alpha^{10})(x^2 + \alpha^{10} x + 1)$$

$$k=9$$

Ex: $q=4, n=21$

Find m s.t. $21 \mid 4^m - 1$ $m=3$ ✓

$\beta \in GF(64)$, primitive

$$\alpha = \beta^3 \Rightarrow \text{ord}(\alpha) = n = 21$$

$$\{0\}, \{1, 4, 16\}, \{2, 8, 11\}, \{3, 12, 6\}, \{5, 20, 17\},$$

$$\{7\}, \{9, 15, 18\}, \{10, 19, 13\}, \{14\}$$

$x + \beta^{21}$ $x + \beta^{42}$

BCH codes over $GF(4)$, $n=21$

$$t=1 \quad g(x) = M_d(x) M_{d^2}(x) \quad \text{deg} = 6, k = 15$$

$$t=2 \quad g(x) = M_d(x) M_{d^2}(x) M_{d^3}(x) \quad k = 12$$

$$t=3 \quad g(x) = M_d M_{d^2} M_{d^3} M_{d^5} \quad k = 9$$

Reed-Solomon Codes

q : prime power $n = q - 1$

$\alpha \in GF(q)$, primitive

$$\text{ord}(\alpha) = n$$

$$\{0\}, \{1\}, \{2\}, \dots, \{q-2\}$$
$$x^{q-1} - 1 = (x - \alpha^0)(x - \alpha^1)(x - \alpha^2) \dots (x - \alpha^{q-2})$$

$i \pmod{q} \text{ mod } (q-1) = \underline{i}$

$$t=1 \quad g(x) = (x - \alpha)(x - \alpha^2) \quad k = n - 2t \quad \text{MDS}$$

$$t \quad \dots \quad g(x) = (x - \alpha)(x - \alpha^2) \dots (x - \alpha^{2t})$$

Polynomial Evaluation Construction of Reed-Solomon Codes

n : blocklength

$\alpha \in GF(q)$, $\text{ord}(\alpha) \geq n$

$$f(x) = f_0 + f_1 x + \dots + f_{k-1} x^{k-1}$$

message = $[f_0 f_1 \dots f_{k-1}]$ $f_i \in GF(q)$ $\underbrace{k}_{\text{length}} = \underline{\underline{q^k}}$

$$\text{Codeword} = [f(\alpha) f(\alpha^2) f(\alpha^3) \dots f(\alpha^n)]$$

length = n

$$\Rightarrow \text{wt}(\text{codeword}) \geq n - (k-1) = n - k + 1$$