

Reed-Solomon Codes

Note Title

n : block length, designed error-correcting capability = t

$$\beta \in GF(2^m), \text{ord}(\beta) \geq n$$

$$H = \begin{bmatrix} 1 & \beta & \beta^2 & \dots & \beta^{n-1} \\ 1 & \beta^2 & (\beta^2)^2 & \dots & (\beta^2)^{n-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \beta^{2t} & (\beta^{2t})^2 & \dots & (\beta^{2t})^{n-1} \end{bmatrix} \quad \text{rank}(H) = 2t$$

$$\text{RS Code} = \left\{ \underline{c} \in \underline{GF(2^m)}^n : H \underline{c}^T = \underline{0} \right\}$$

$$H \underline{c}^T = \underline{0}$$

first row:

$$\underline{c} = [c_0 \ c_1 \ c_2 \ \dots \ c_{n-1}]$$
$$c(x) = c_0 + c_1 x + \dots + c_{n-1} x^{n-1} \quad c_i \in GF(2^m)$$

$$2^{\text{nd}} \text{ row: } c(\beta^2) = 0 \Leftrightarrow (x + \beta^2) \mid c(x)$$

\vdots

$$2t\text{-th row: } c(\beta^{2^t}) = 0 \Leftrightarrow (x + \beta^{2^t}) \mid c(x)$$

$$\Leftrightarrow \underbrace{(x + \beta)(x + \beta^2) \dots (x + \beta^{2^t})}_{g(x)} \mid c(x)$$

deg = $2t$

$g(x)$: generator polynomial

$$\text{RS code} = \left\{ c(x) = m(x)g(x) : \right.$$

$$m(x) = m_0 + m_1x + \dots + m_{k-1}x^{k-1}$$

$$\left. \begin{array}{l} k = n - 2t \end{array} \right\}$$

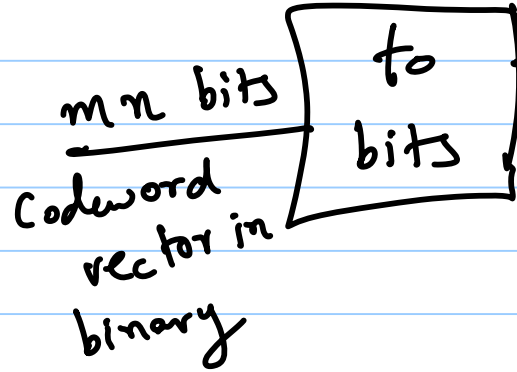
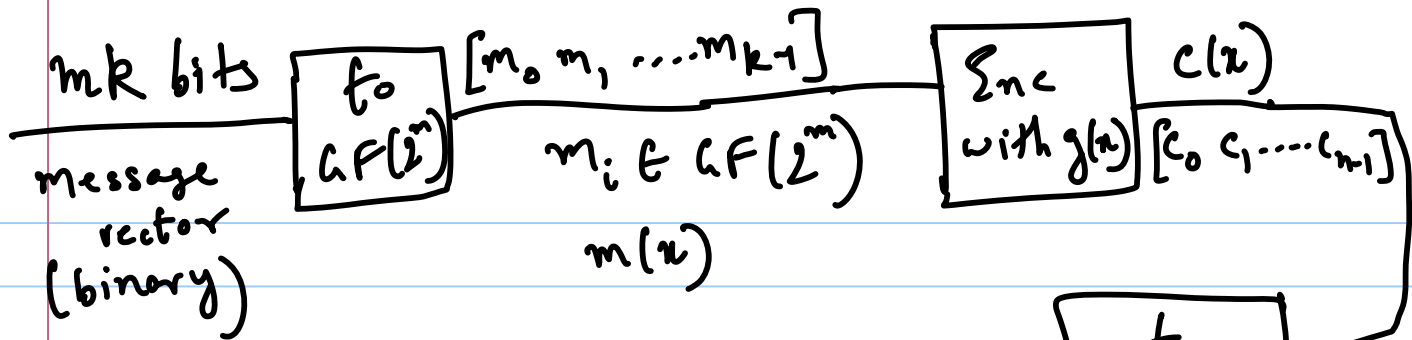
$$\text{Dimension} = n - 2t = k$$

$$\text{Minimum distance} = 2t + 1 = n - k + 1$$

\Rightarrow RS codes have parameters $(n, k, d = n - k + 1)$

\downarrow \downarrow \downarrow

MDS codes n : elements of $\text{GF}(2^m)$

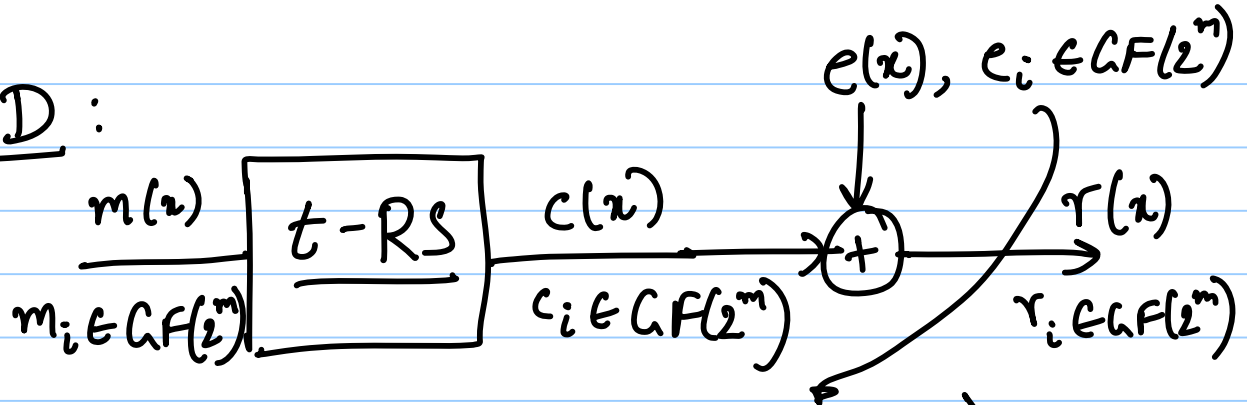


$m=8, GF(256)$
 Ex: $k=239, n=255$

$mk = 1912$ bits, $nm = 2040$ bits

$n = \text{ord}(\beta) \Rightarrow$ RS codes will be cyclic.

BDD:



w errors $\Rightarrow \text{wt}([e_0 \ e_1 \ \dots \ e_{n-1}]) = w$

$$w \leq t \quad P_r = \binom{n}{w} p_s^w (1-p_s)^{n-w}$$

$$P_r(e_0 \neq 0) = 1 - (1-p)^m = \underline{\underline{p_s}} \quad \text{prob of symbol error}$$

Form of $e(x)$: $e(x) = E_1 x^{i_1} + E_2 x^{i_2} + \dots + E_w x^{i_w}$
 with w errors
 $E_i \in GF(2^m)$

Decoding task: Find i_1, i_2, \dots, i_w and
 E_1, E_2, \dots, E_w .

Syndrome:

$$S_1 = r(\beta) = E_1 \beta^{i_1} + E_2 \beta^{i_2} + \dots + E_w \beta^{i_w}$$

$$S_2 = r(\beta^2) = E_1 \beta^{2i_1} + E_2 \beta^{2i_2} + \dots + E_w \beta^{2i_w}$$

$$\vdots$$

$$S_{2^t} = r(\beta^{2^t}) = E_1 \beta^{2^t i_1} + \dots + E_w \beta^{2^t i_w}$$

Error locator: $X_j = \beta^{ij}$
 $j=1, 2, \dots, w$

Error values: E_j

Known

$$\begin{aligned} S_1 &= E_1 X_1 + E_2 X_2 + \dots + E_w X_w \\ S_2 &= E_1 X_1^2 + E_2 X_2^2 + \dots + E_w X_w^2 \\ &\vdots \\ S_{2t} &= E_1 X_1^{2t} + E_2 X_2^{2t} + \dots + E_w X_w^{2t} \end{aligned}$$

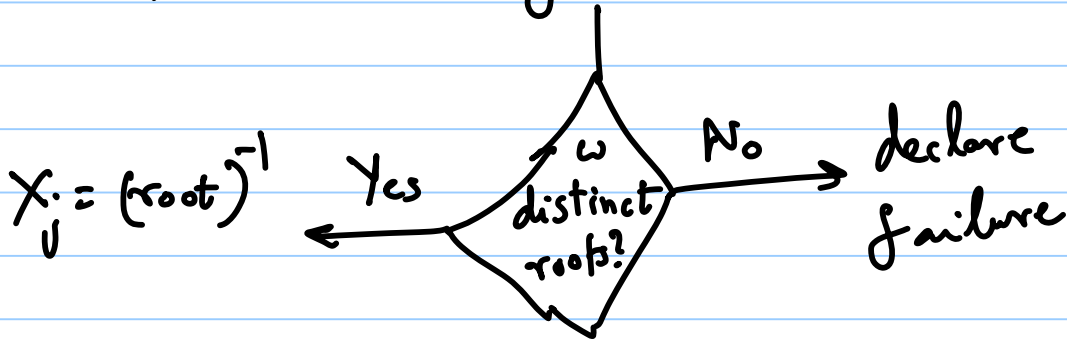
Error locator poly, $\sigma(x) = (1 + X_1 x)(1 + X_2 x) \dots (1 + X_w x)$
 $(+ \sigma_1 x + \dots + \sigma_w x^w)$

$$S(x) = S_1 x + S_2 x^2 + \dots + S_{2t} x^{2t}$$

$S(x) \sigma(x)$: coeffs of $x^{\omega+1}, \dots, x^{2t}$ are zero

Find $\sigma_1, \sigma_2, \dots, \sigma_\omega$.

Find roots of $\sigma(x)$ in $GF(2^m)$.



→ Use w Syndrome equations (now linear in E_j)

to find E_1, E_2, \dots, E_w .

